

# Evaluating I/O Acceleration Mechanisms of SX-Aurora TSUBASA

Yuta Sasaki\*, Ayumu Ishizuka\*, Mulya Agung<sup>‡</sup>, Hiroyuki Takizawa<sup>‡</sup>\*

- \* Graduate School of Information Sciences, Tohoku University, {ysasaki,ishizuka}@hpc.is.tohoku.ac.jp
- ‡ Cyberscience Center, Tohoku University, agung@hpc.is.tohoku.ac.jp, takizawa@tohoku.ac.jp

### Outline

- **■**Introduction
- **■**Target System
- **■**Performance Evaluation
- ■Use Case of I/O Acceleration
- **■**Conclusions and Future Work

### Introduction

### Heterogeneous HPC systems

- Promising to increase their computational performance
- Bring a new challenge in achieving high file I/O performance

### ■I/O operations are processed by the collaboration among...

- Host processor is responsible for handling system calls
- Some processors **not fully supporting OS functions** invoked via system calls
  - Ask the host to manage system calls such as file I/O operations

### Motivation

- ■File I/O performance could be a performance bottleneck
  - Intermediate simulation results are periodically stored into files in practical numerical simulations
- Auto-tuning is required for I/O performance as well as computational performance
  - Understanding the I/O characteristics of heterogeneous systems is needed

- **■**Introduction
- **■**Target System
- ■Performance Evaluation
- ■Use Case of I/O Acceleration
- ■Conclusions and Future Work

## Target System

#### ■NEC SX-Aurora TSUBASA (SX-AT)

- A latest vector computing system with heterogeneous configuration
- Vector hosts (VHs): x86 host processor
- Vector engines (VEs): NEC's proprietary vector processor
  - General-purpose processor
    - Programmers can execute the whole application code
  - Rely on the VH to provide OS functionality
    - Need additional PCIe data transfer between VH and VE for I/O operations
    - Potentially cause non-negligible overhead



NEC SX-Aurora TSUBASA Vector Engine\*

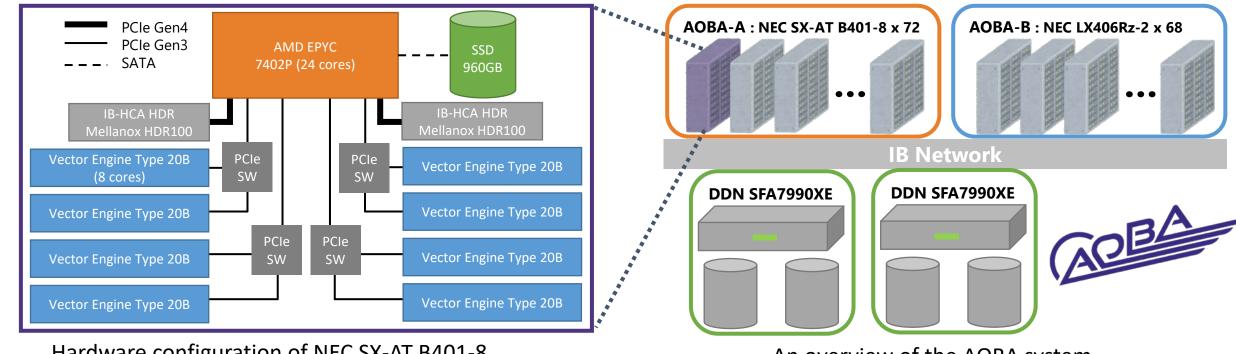
### ■Two I/O acceleration mechanisms

- Accelerated I/O (AccIO)<sup>[3]</sup>
- ScaTeFS VE Direct I/O (DirIO)<sup>[3]</sup>
- Implemented as libraries to transparently replace file I/O system calls
  - Without any modification of application code

# **System Configuration**

### ■AOBA system installed at the Cyberscience Center, Tohoku University

- The remote storage system is relatively small of only two DDN SFA7990XE
- Local SSD is also equipped on each node

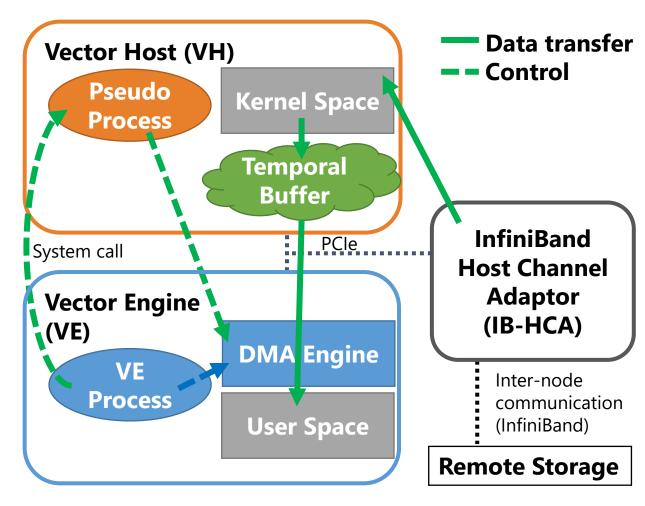


Hardware configuration of NEC SX-AT B401-8

An overview of the AOBA system

# Normal I/O (without acceleration)

- Most system calls on VEs are offloaded to user space daemons running on VH (pseudo process)
  - VE needs additional PCIe data transfer to access a file
- Data transfers for file read operations
  - Read from the file to a kernel space buffer in VH
  - Copied to a temporal buffer allocated in the VH user space
  - Sent from the VH to the VE via PCle interconnect

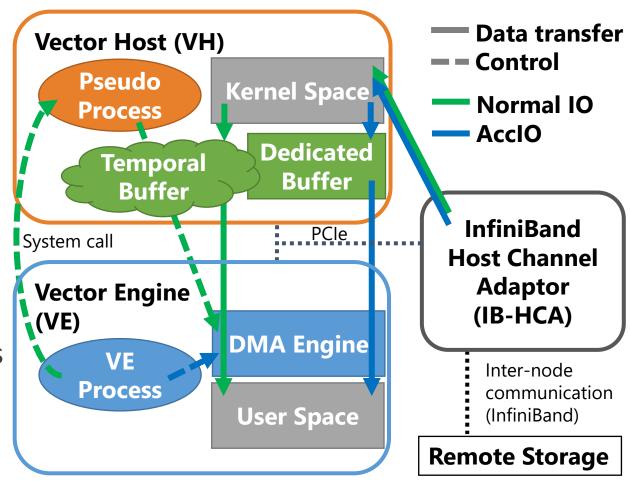


Data read from the remote storage via IB-HCA

# Accelerated I/O (AccIO)

### Acceleration through elimination of two overheads

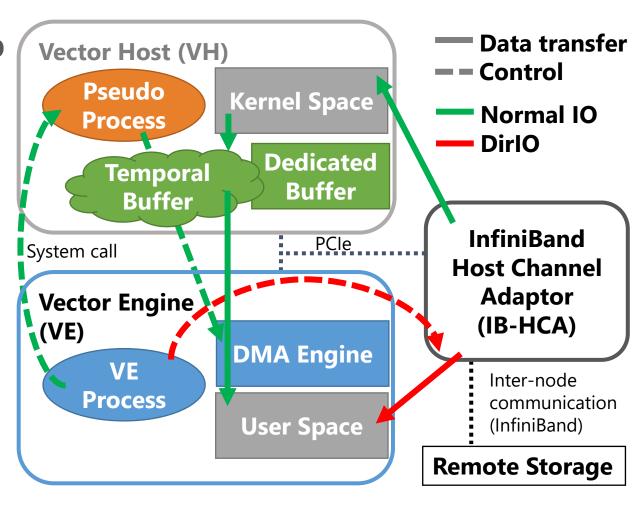
- Pseudo process in VH controls VE's direct memory access (DMA) engine
  - AccIO allows a user process running on VE (VE process) to directly control the DMA engine
  - Improve the efficiency of data transfers between a VH and a VE via PCIe
- A temporal buffer in the VH user space is allocated and released whenever a read system call is invoked
  - Page-locked and dedicated buffer is allocated once and reused



Data read from the remote storage via IB-HCA (AccIO)

# ScaTeFS VE Direct I/O (DirIO)

- Acceleration through direct access to ScaTeFS parallel file system
  - VE process is capable of directly communicating with IB-HCA
    - Without offloading relevant system calls to VH
    - When the I/O size is greater than 1MB
- ■AccIO and DirIO cannot be used at the same time
  - Need to decide which one should be enabled for a given application



Data read from the remote storage via IB-HCA (DirIO)

- Introduction
- ■Target System
- **■**Performance Evaluation
- ■Use Case of I/O Acceleration
- ■Conclusions and Future Work

### **Evaluation**

### **■**Evaluate the file I/O performance of the ScaTeFS parallel file system<sup>[4]</sup>

- Performance values shown in this work depend on the system configuration
  - i.e. Theoretical peak I/O bandwidth of ScaTeFS scales with the number of I/O servers
- We do not intend to make comparisons with other systems

### ■How does the I/O bandwidth change with the I/O acceleration mechanisms and application behaviors?

- Three I/O modes (Normal I/O, AccIO, DirIO)
- IOR benchmark<sup>[1]</sup>
  - Widely used to discuss the I/O performance of HPC systems
  - Reproduce the I/O behaviors of various applications by adjusting parameters

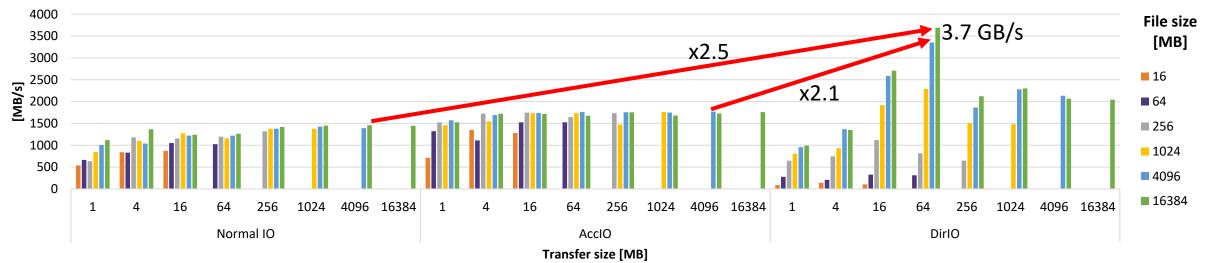
iWAPT2021

• File size, Transfer size and Single-shared-file/File-per-process etc...

# Single-process Write Performance

### ■Performance change by the file size f and the transfer size t ( $f \ge t$ )

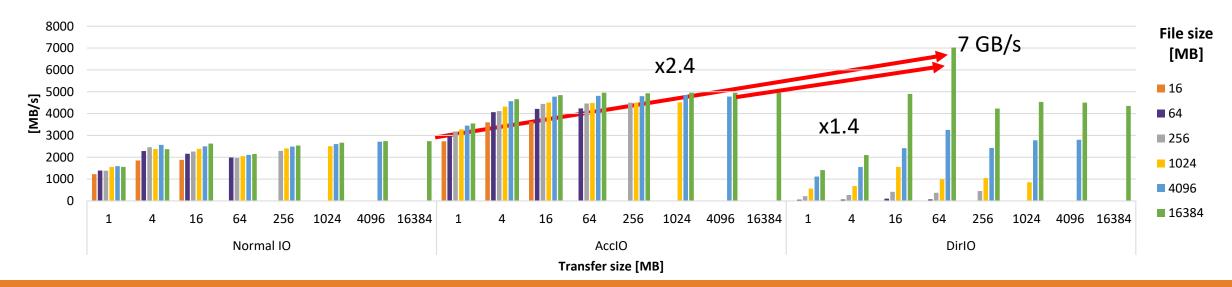
- AcclO performance is higher than the normal I/O performance
  - AccIO improves the efficiency of the data copy between VHs and VEs
- DirlO performance is much more sensitive to file size f and transfer size t
  - f < 1 GB : DirlO performance is worse even than the normal I/O
  - Significantly increases with f and t and reaches 3.7 GB (t = 64MB, f = 16GB)
  - Outperforms the others when writing a sufficiently large file with appropriate t



# Single-process Read Performance

### ■The read performance characteristics are similar to those for write

- DirlO needs an even larger file size of 16 Gbytes to outperform AcclO
  - For applications that need to frequently read small files of less than 16GB, it is worth examining to disable DirlO in order to achieve high sustained performance.
- DirlO can achieve the best performance with t = 64 MB
  - ScaTeFS stripe size: 4 MB x num of parallel I/O operations for a single process: 16 = 64 MB



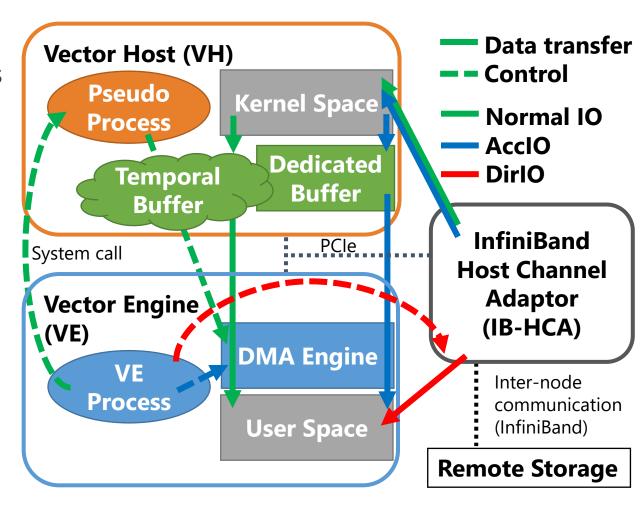
# Discussion / Why is DirlO sensitive to t?

### ■System call handling is not vectorizable

- VH is suited for non-vectorizable operations
- VE's overhead would become larger

# ■ DirIO enables VE to communicate directly with IB-HCA without offloading

- When the transfer size is small...
  - Overhead is non-negligible and rather dominant in the I/O time
  - If DirIO is disabled, VH does caching and prefetching I/O data while VE does not support
- As the transfer size increases...
  - VE's overhead becomes relatively smaller
  - the benefit of direct communication between VE and IB-HCA outweighs the overhead



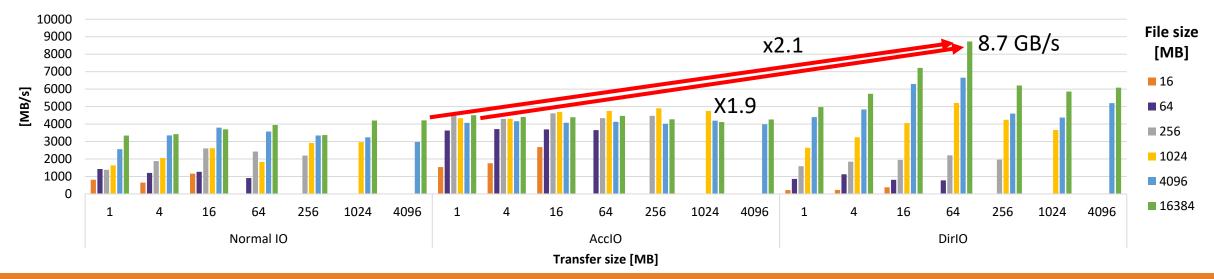
# Multi-process Write Performance (1VE)

### **■**Discuss the I/O performance of an MPI application

- All of the 8 cores on a VE are used to access files in parallel
- Each process writes a different file (file-per-process mode)
- File size indicates the size of each file (not the total size)

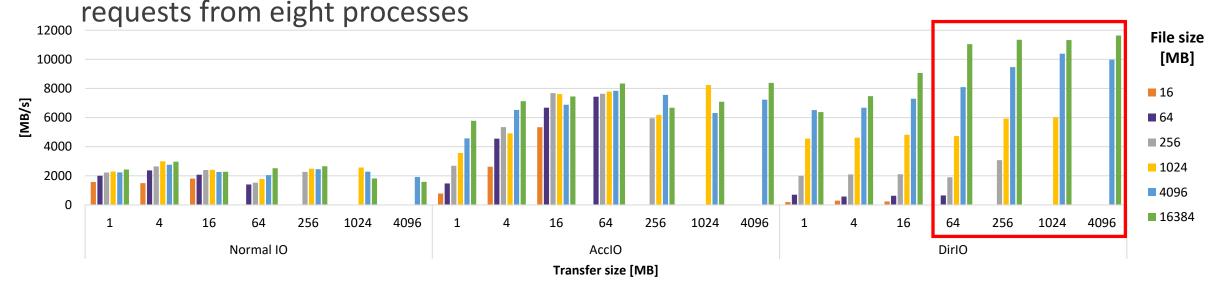
### ■Aggregated bandwidth of 8 processes

• Reaches about 8.7 GB/s at t = 64MB and f = 16GB as with the single-process



# Multi-process Read Performance (1VE)

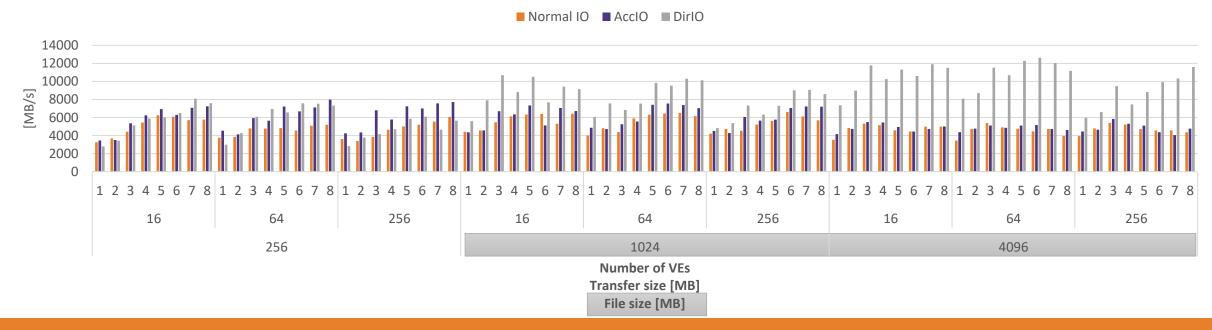
- ■No remarkable peak at the transfer size of 64 Mbytes
  - Performance remains high for a larger transfer size
- ■Sustained bandwidth of about 12 GB/s when reading the largest file
  - The theoretical peak bandwidth of PCIe Gen3 is 16 GB/s
  - The interconnect bandwidth is almost saturated by a lot of concurrent file access



# Write Performance with Multiple VEs

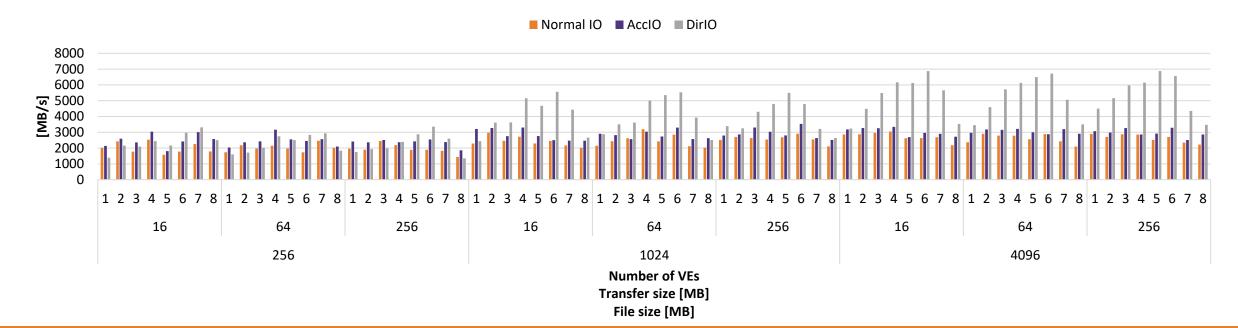
### ■When increasing the number of VEs (each VE executes 8 processes)

- DirlO shows the best performance for the file size of 1 GB or larger
- When the file size and the number of MPI processes are large,
   DirlO shows outstanding performance compared to AcclO.



# Read Performance with Multiple VEs

- ■When the number of VEs is large,
  DirlO shows the best performance even for a small file of 1 GB
- ■The impacts of skipping data copy between VH and VEs become more significant when using multiple VEs



# Single-shared-file Write Performance

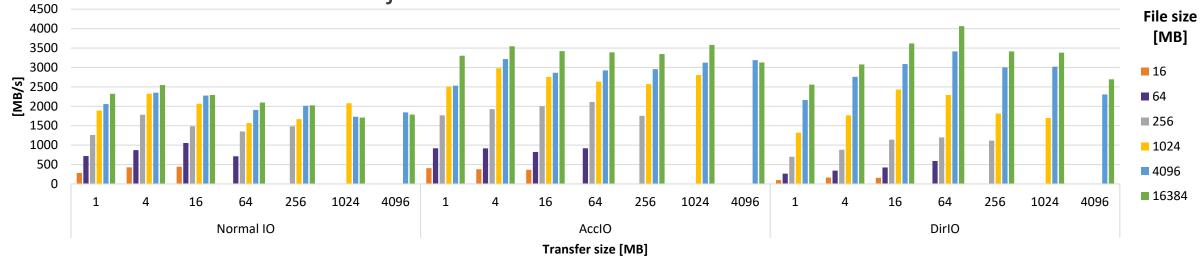
- ■All 8 processes share a single file ("single-shared-file" mode)
  - File size represents the size that each process writes
- ■AccIO performance is lower than that in the "file-per-process" mode
  - May be due to file blocking
  - The performance degradation of DirlO is small



# Single-shared-file Read Performance

- ■Unlike the write performance, the performance of AccIO and DirIO become lower than those in the file-per-process mode
- ■DirIO is likely to be effective when a lot of concurrent I/O operations are executed simultaneously, even if the size of each file is small

• This property would be necessary for the operation of HPC systems because a lot of various jobs run at the same time.



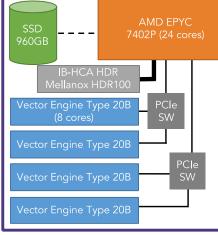
# Write Performance to Local SSD

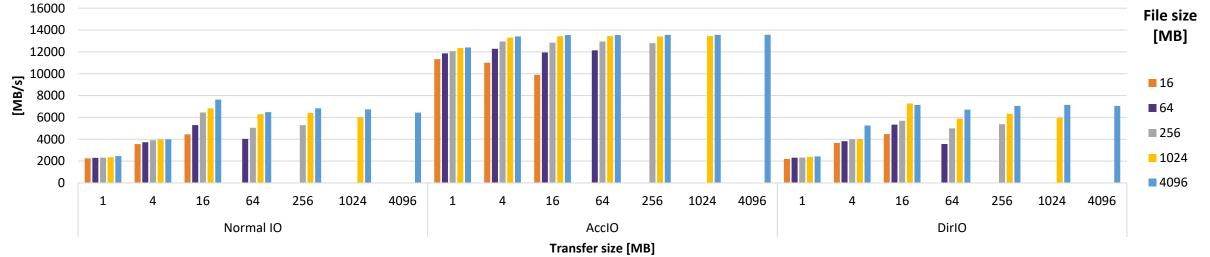
### ■Each node of AOBA is equipped with a local SSD

All VEs in the node share and access the SSD via the PCIe

### **■**Only AccIO accelerates write performance

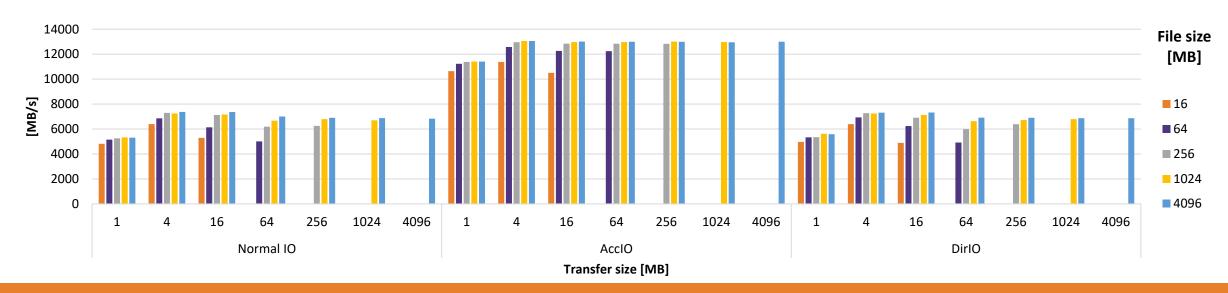
- AccIO can improve the data transfer performance between a VH and a VE via via the PCIe interconnect
- DirlO does not significantly affect the I/O performance because the access to local storage does not go through the IB-HCA





# Read Performance from Local SSD

- ■As with write, only AccIO accelerates read performance
- ■If an application accesses local storages more intensively than the parallel file system
  - AccIO should be used to achieve higher I/O performance



- Introduction
- ■Target System
- ■Performance Evaluation
- ■Use Case of I/O Acceleration
- **■**Conclusions and Future Work

# Use Case of I/O Acceleration

■Use a real-world MPI application of flood simulation<sup>[9] [10]</sup> to discuss the performance benefit of I/O acceleration in practical use

iWAPT2021

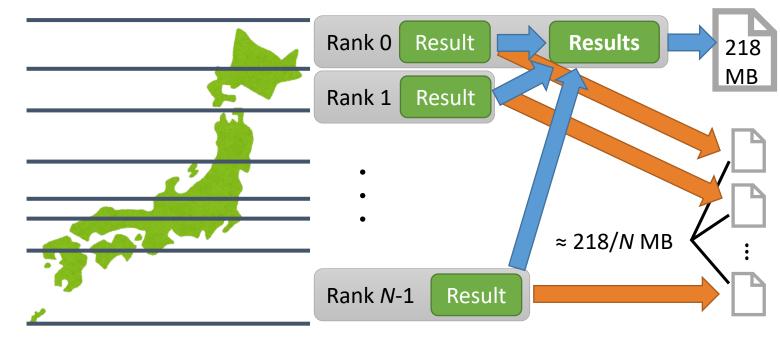
- Divide the land areas into sections
- Predict flood damage in parallel with MPI

### Gathered mode (Original)

 Gather intermediate results to Rank 0 and writes the gathered results to a file

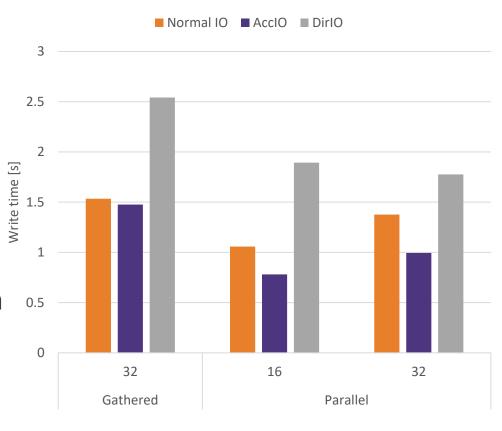
### Parallel mode

 Write partial results to a different file



### Discussion

- ■Write time of the three I/O modes with 16 or 32 MPI processes
  - In parallel mode, the write time is the total time spent on writing all files
- ■Users may consider that DirIO can achieve higher I/O performance in any cases
  - DirIO could degrade the I/O performance for accessing small files as suggested by the evaluation results with the IOR benchmark
- ■Need to carefully select either AccIO or DirIO and its parameters
  - Considering the file access behaviors



Write performance in flood simulation

### Conclusions and Future Works

- ■The first investigation into effects of AccIO and DirIO on I/O of SX-AT
  - Discussed proper use of AccIO and DirIO for a real-world application
- ■Our evaluation results clearly show that the two I/O acceleration mechanisms have their own pros and cons
  - Appropriately used considering the application behaviors and system configuration.
  - Clarified the demand for auto-tuning technology to appropriately select either of the two I/O acceleration mechanisms of SX-AT
  - Abstraction and auto-tuning of those mechanisms will be discussed in our future work

# Acknowledgments

- ■This work is partially supported by
  - MEXT Next Generation High-Performance Computing Infrastructures and Applications R&D Program "R&D of A Quantum-Annealing-Assisted Next Generation HPC Infrastructure and its Applications"
  - Grant-in-Aid for Scientific Research(A) #20H00593
- ■The authors would like to thank Takashi Sato and Yuta Watanabe of NEC corporation for their technical supports and fruitful discussions